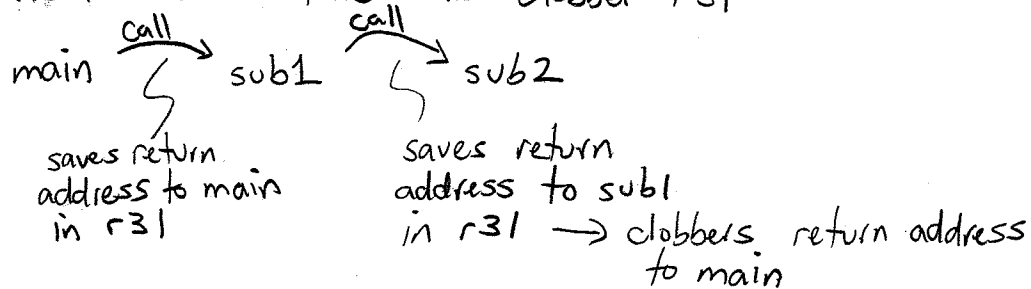
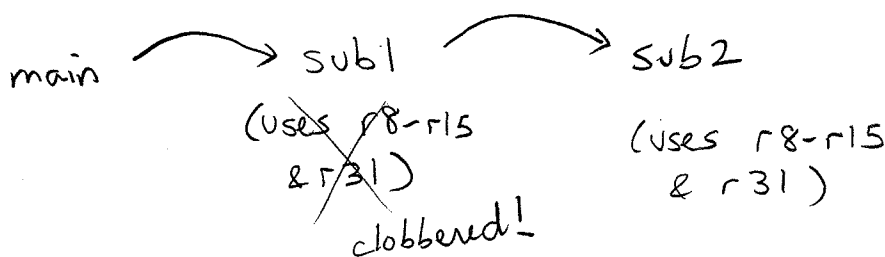


## Stacks and Subroutines

Problem: nested subroutine calls clobber r31



other registers (used in subroutine) get clobbered too!



Solution: do two things

- save important registers on stack before subroutine call
- restore registers from stack after subroutine returns

Stack: ideally, an unbounded region of memory where we store values in last-in-first-out (LIFO) order

practically, there will be a memory limit

no preset size limit

reversing

eg: long driveway full of cars, last one in is first one out!

convention: r27 is "stack pointer" or sp used to keep track of top of stack

# Example

```
.equ DRAM_END, 0x0080 0000
```

```

_start: movia sp, SRAM_END  ← init stack
        }
        call sub1
        }
        call sub3
        }
stop:   br stop

```

```

sub1:   }
        } ← save r31
        } ① addi sp, sp, -4
        } stw r31, 0(sp)
        call sub2
        }
        } ← ret

```

```

sub2:   }
        } ← restore r31
        } ② ldw r31, 0(sp)
        } addi sp, sp, 4
        }

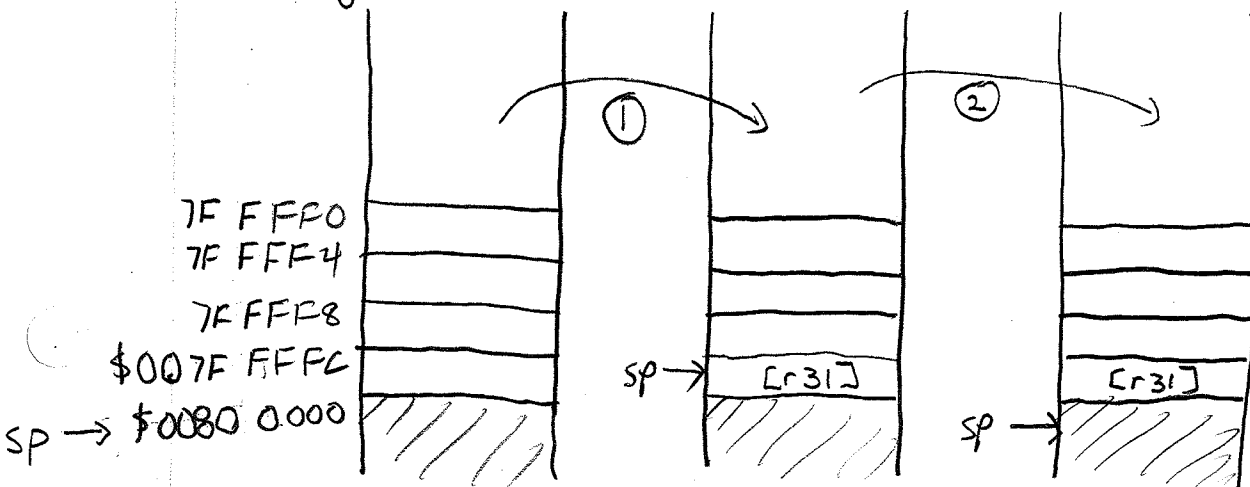
```

```

sub3:   }
        } ← ret

```

Memory:



convention	r2, r3	return value
	r4 - r7	incoming params
	<u>r8 - r15</u>	caller-saved
	subr.	∴ safe to use inside subroutine
caller    employer		
↓        ↓		
callee   employer	<u>r16 - r23</u>	callee-saved
	main	∴ subroutine must save + restore if it changes them
	r1, r24 - r31	don't use (mostly specialized like r27 = sp, r31 = ra, ...)

Saving multiple registers

```

sub1: addi sp, sp, -4
      stw r31, 0(sp)

```

} save ret-addr (sub1 will call sub2)

} sub1 needs r4, r8, r9

Can also save r16-r23 if needed by sub1 (callee saved)

caller-saved →

```

addi sp, sp, -12
stw r4, 8(sp)
stw r8, 4(sp)
stw r9, 0(sp)

```

} save regs needed by sub1

sub2 clobbers (r4, r8, ..., r15, r31) result is in r2, r3 →

```

mov r4, someParamForSub2
call sub2

```

} call sub2

caller-saved →

```

ldw r9, 0(sp)
ldw r8, 4(sp)
ldw r4, 8(sp)
addi sp, sp, 12

```

} restore regs for sub1 in reverse order

} examine sub2 result in r2, r3

```

ldw r31, 0(sp)
addi sp, sp, 4
ret

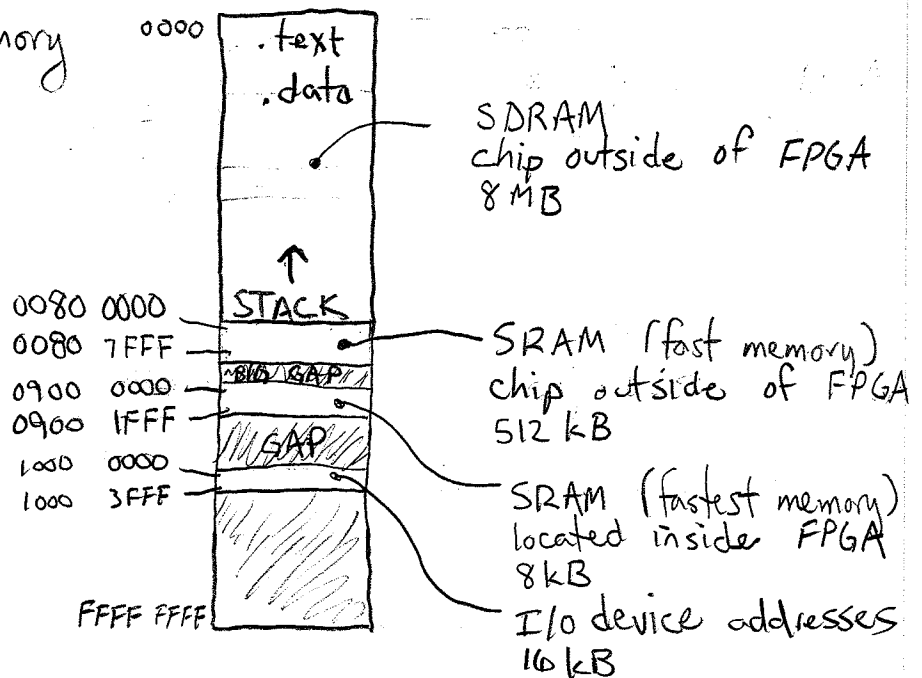
```

} restore ret-addr & return from sub1

(callee saved) also restore r16-r23 if used by sub1



## Bigger view of memory

stack rules

- ① sp always points to top of stack - most recent item put in stack

eg1: `ldw r2, 0(sp)` — always checks top of stack

eg2: `addi sp, sp, -4`  
`stw r31, 0(sp)` — places new item on top of stack

- ② all stack operations are done manually with ordinary instructions - no magic "stack instructions" or "hardware stack" to do things in hidden/automatic way for Nios II
- ③ initially, sp should point to just past the end of stack memory
- ④ as data { added to stack, { sp gets smaller PUSH  
{ removed from stack, { sp gets bigger PULL
- ⑤ MAINTAIN BALANCE always PULL what you PUSH in equal amounts & in reverse order